



Let $p_{0,i}$ ($i=1, 2, \dots, dc-1$) be the probability of 0 of i th input message.
 p_0 be the probability of 0 of output message.

Note that: $[p_0 \ p_1] = [p_{0,1} \ p_{1,1}] \otimes [p_{0,2} \ p_{1,2}] \otimes \dots \otimes [p_{0,dc-1} \ p_{1,dc-1}]$,
 where \otimes denotes circular convolution.

Hence:

$$\mathcal{F}[p_0, p_1] = \sum_{i=1}^{dc-1} \mathcal{F}[p_{0,i}, p_{1,i}] \quad (1)$$

where \mathcal{F} is DFT, and multiplication here is entrywise.

Recall DFT:

$$X_k = \sum_{n=0}^{N-1} x_n e^{-j\frac{\pi}{N} kn}$$

Hence, the 2-point DFT of $[p_0, p_1]$ is:

$$[p_0 + p_1, p_0 + p_1 e^{-j\frac{\pi}{2}}] = [p_0 + p_1, p_0 - p_1]$$

Equation (1) gives:

$$[p_0 + p_1, p_0 - p_1] = [\underbrace{\sum_{i=1}^{dc-1} (p_{0,i} + p_{1,i})}_{=1}, \underbrace{\sum_{i=1}^{dc-1} (p_{0,i} - p_{1,i})}_{=1}]$$

$$\therefore p_0 - p_1 = \sum_{i=1}^{dc-1} (p_{0,i} - p_{1,i})$$

$$\text{Let } l_0 = (\log \frac{p_0}{p_1}), \text{ then: } \frac{p_0}{p_1} = e^{l_0} \Rightarrow p_0 = p_1 e^{l_0} \Rightarrow \begin{cases} p_0 - p_1 = p_1 (e^{l_0} - 1) \\ p_0 + p_1 = p_1 (e^{l_0} + 1) \end{cases}$$

$$\therefore p_0 - p_1 = \frac{p_0 - p_1}{p_0 + p_1} = \frac{e^{l_0} - 1}{e^{l_0} + 1} = \tanh(\frac{l_0}{2}) \quad [\tanh(x) = \frac{e^{2x} - 1}{e^{2x} + 1}]$$

$$\therefore \tanh(\frac{l_0}{2}) = \sum_{i=1}^{dc-1} \tanh(\frac{l_0}{2}) \Rightarrow l_0 = 2 \cdot \tanh^{-1} \left(\sum_{i=1}^{dc-1} \tanh(\frac{l_0}{2}) \right) \quad \square$$